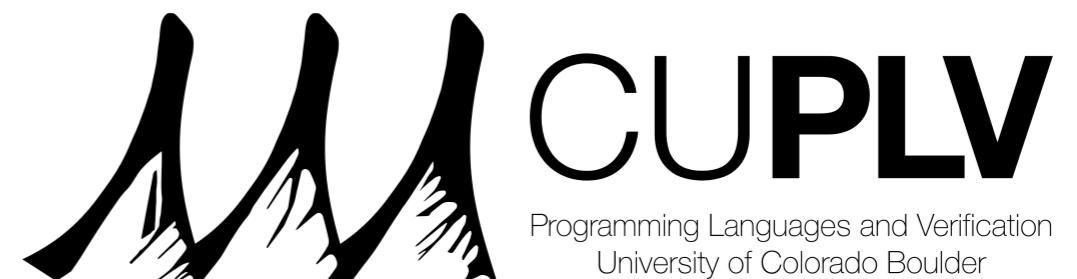


Static Analysis with Demand-Driven Value Refinement

Benno Stein, **Benjamin Barslev Nielsen**, Bor-Yuh Evan Chang & Anders Møller



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Sound static analysis for JavaScript

- Static analysis for JavaScript is very challenging

o[m]()

Sound static analysis for JavaScript

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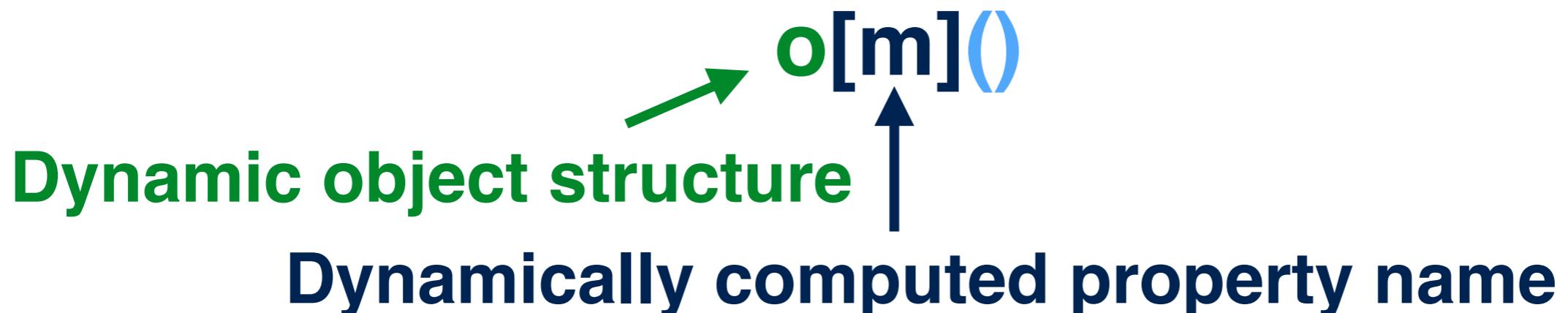


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Dynamic object structure

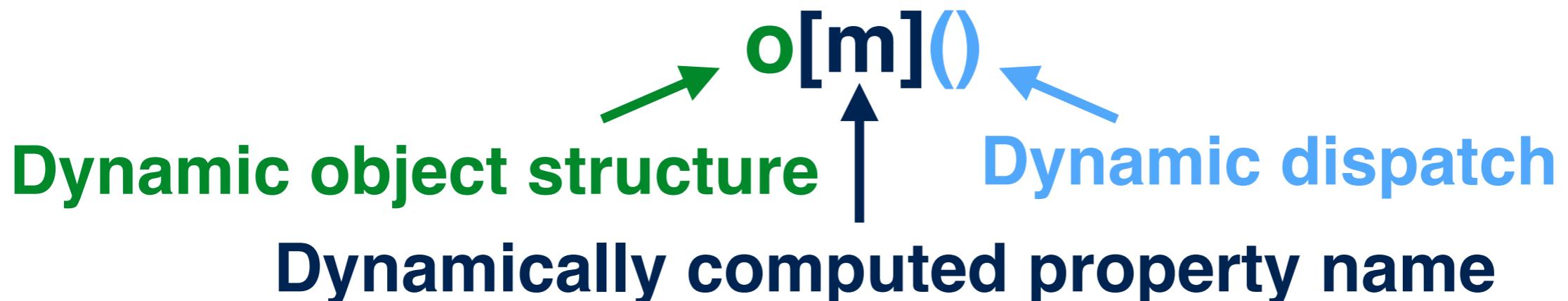
Sound static analysis for JavaScript

- Static analysis for JavaScript is very challenging



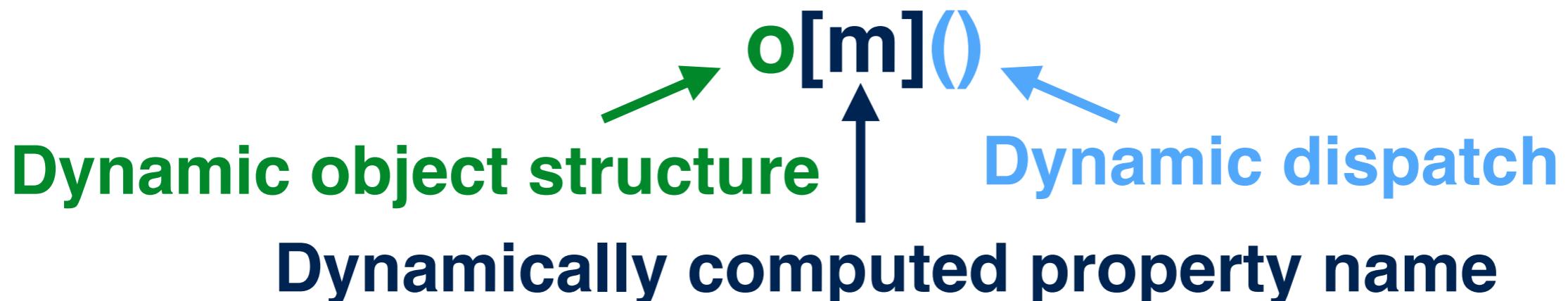
Sound static analysis for JavaScript

- Static analysis for JavaScript is very challenging



Sound static analysis for JavaScript

- Static analysis for JavaScript is very challenging



- Critical precision losses renders analysis useless
 - Too much spurious data-flow

State-of-the-art data-flow analyzers

- Fail to analyze load of some very popular libraries
 - Critical precision losses occur
- Common characteristics
 - Forwards whole program analysis
 - Tracks data-flow, e.g., strings, functions and other objects
 - Non-relational
 - Aims to mitigate critical precision losses by:
 - Context sensitivity
 - Syntactic patterns and special-case techniques

Critical code example

Example program

Analysis state

$\text{func} = \text{o1}[\text{name}]$

$\text{o2}[\text{name}] = \text{func}$

$\text{o2}.\text{foo}(\dots)$

Critical code example

Analysis state

```
o1 = {foo: f1, bar: f2}
```

```
name = Tstr
```

```
o2 = {}
```

Example program

```
func = o1[name]
```

```
o2[name] = func
```

```
o2.foo(...)
```

Critical code example

Example program

Analysis state

```
o1 = {foo: f1, bar: f2}
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```
name = Tstr
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```
o2 = {}
```

```
func = f1|f2
```

```
func = o1[name]
```

```
o2[name] = func
```

```
o2.foo(...)
```

Critical code example

Example program

Analysis state

```
o1 = {foo: f1, bar: f2}
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name = Tstr
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```
o2 = { Tstr : f1|f2}
```

```
func = f1|f2
```

```
func = o1[name]
```

```
o2[name] = func
```



```
o2.foo(...)
```

Critical code example

Example program

Analysis state

```
o1 = {foo: f1, bar: f2}
```

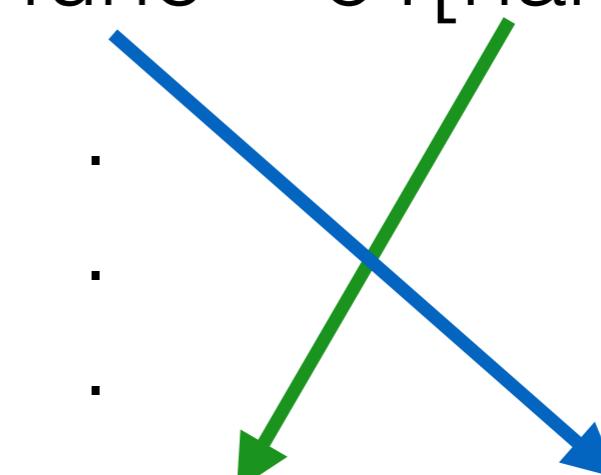
```
name = Tstr
```

```
o2 = { Tstr : f1|f2}
```

```
func = f1|f2
```

```
func = o1[name]
```

```
o2[name] = func
```

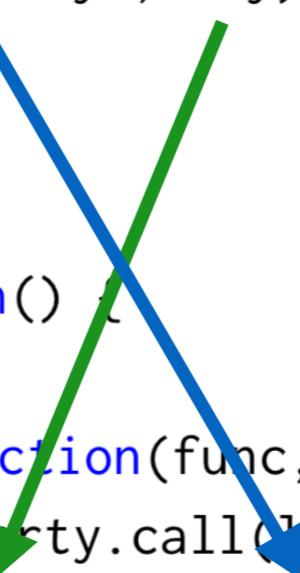


Resolves both f_1 and f_2

→ `o2.foo(...)`

The Lodash library

```
1 function baseFor(object, iteratee) {  
2     ...  
3     while (length--) {  
4         var key = props[++index];  
5         iteratee(object[key], key)  
6     }  
7 }  
8  
9 mixin(lodash, (function() {  
10     var source = {};  
11     baseFor(lodash, function(func, methodName) {  
12         if (!hasOwnProperty.call(lodash.prototype, methodName)) {  
13             source[methodName] = func;  
14         }  
15     });  
16     return source;  
17 }()));
```



Critical code example

Example program

Analysis state

```
o1 = {foo: f1, bar: f2}
```

```
name = Tstr
```

```
o2 = {}
```

```
func = f1|f2
```



```
func = o1[name]
```

```
o2[name] = func
```

```
o2.foo(...)
```

Critical code example

Example program

Analysis state

```
o1 = {foo: f1, bar: f2}
```

```
name = Tstr
```

```
o2 = {}
```

```
func = f1|f2
```

```
func = o1[name]
```

```
o2[name] = func
```

```
.
```

```
.
```

```
.
```

```
o2.foo(...)
```

Demand-driven value refinement

Regain relational information through refinement queries

Without modifying base analysis domain

Refinement query: What is x , when $y \mapsto \hat{v}$?

What value can variable x have,
given that y has value \hat{v} ?

Critical code example

Analysis state

```
o1 = {foo: f1, bar: f2}  
name = Tstr  
o2 = {}  
func = f1|f2
```

Example program

```
func = o1[name]  
.  
.  
.  
→ o2[name] = func  
.  
.  
.
```

o2.foo(...)

Critical code example

Analysis state

$o1 = \{foo: f1, bar: f2\}$

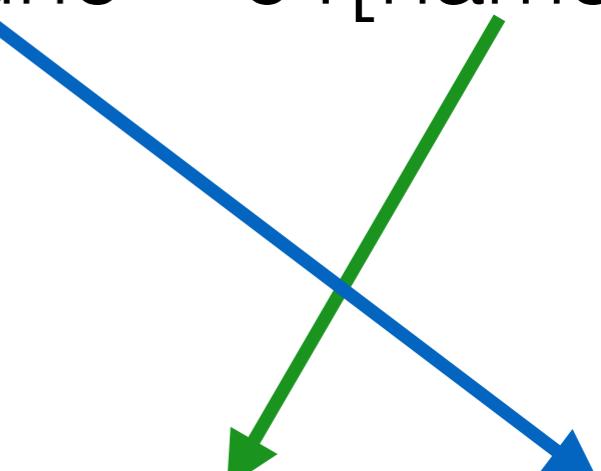
$name = T_{str}$

$o2 = \{\}$

$func = f1 | f2$

Example program

$func = o1[name]$



What is name, when $func \mapsto f1$? → $o2[name] = func$

What is name, when $func \mapsto f2$?

$o2.foo(\dots)$

Backwards abstract interpreter for value refinement

- Backwards goal-directed from the query location
- Separation logic based abstract domain
 - Intuitionistic - constraints hold for all extensions
 - Special symbolic variable RES represents value being refined

symbolic variables $\hat{x}, \hat{y}, \hat{z}, \text{RES} \in \widehat{\text{Var}}$

symbolic stores $\varphi \in \widehat{\text{Store}} ::= \hat{h} \wedge \pi \mid \varphi_1 \vee \varphi_2$

heap constraints $\hat{h} ::= \text{true} \mid \text{unalloc}(\hat{x}) \mid x \mapsto \hat{x} \mid \hat{x}_1[\hat{x}_2] \mapsto \hat{x}_3 \mid \hat{h}_1 * \hat{h}_2$

pure constraints $\pi ::= \text{true} \mid \hat{e} \mid \pi_1 \wedge \pi_2$

symbolic expressions $\hat{e} \in \widehat{\text{Expr}} ::= \hat{x} \mid \hat{v} \mid \hat{e}_1 \oplus \hat{e}_2$

Backwards abstract interpreter for value refinement

- Based on refutation sound Hoare triples $\langle \varphi \rangle s \langle \varphi' \rangle$
- Refutation soundness:
 - For all concrete runs where φ' holds after s , the state before s must satisfy φ .
- Encoding refinement queries:

What is x , when $y \mapsto \hat{v}$? $\rightsquigarrow \langle x \mapsto \text{RES} * y \mapsto \hat{y} \wedge \hat{y} = \hat{v} \rangle$

Critical code example

`func = o1[name]`

→ `o2[name] = func`

Critical code example

Refinement query: What is name, when $\text{func} \mapsto \text{f1}$?

$\text{func} = \text{o1}[\text{name}]$

→ $\text{o2}[\text{name}] = \text{func}$

Critical code example

Refinement query: What is name, when $\text{func} \mapsto f1$?

$\text{func} = o1[\text{name}]$

$\langle \text{name} \mapsto \text{RES} * \text{func} \mapsto \widehat{\text{func}} \wedge \widehat{\text{func}} = f1 \rangle$

→ $o2[\text{name}] = \text{func}$

Critical code example

Refinement query: What is name, when $\text{func} \mapsto f_1$?

$$\langle \text{name} \mapsto \text{RES} * \widehat{o_1} \mapsto \widehat{o_1} * \widehat{o_1} [\text{RES}] \mapsto \widehat{\text{func}} \wedge \widehat{\text{func}} = f_1 \rangle$$

$$\text{func} = o_1[\text{name}]$$

$$\langle \text{name} \mapsto \text{RES} * \widehat{\text{func}} \mapsto \widehat{\text{func}} \wedge \widehat{\text{func}} = f_1 \rangle$$

$$\rightarrow o_2[\text{name}] = \text{func}$$

Leveraging forwards analysis state

Analysis state

$o1 = \{foo: f1, bar: f2\}$

$$\langle \text{name} \mapsto \text{RES} * o1 \mapsto \widehat{o1} * \widehat{o1} \mid \text{RES} \mapsto \widehat{\text{func}} \wedge \widehat{\text{func}} = f1 \rangle$$

Refinement result is the values of RES satisfying:

$$o1[\text{RES}] = f1$$

Refinement result: “foo”

Critical code example

Analysis state

```
o1 = {foo: f1, bar: f2}  
name = Tstr  
o2 = {}  
func = f1|f2
```

Example program

```
func = o1[name]  
.  
.  
.  
→ o2[name] = func  
.  
.  
.  
o2.foo(...)
```

Critical code example

Analysis state

$o1 = \{foo: f1, bar: f2\}$

$name = T_{str}$

$o2 = \{\}$

$func = f1 | f2$

What is name, when $func \mapsto f1$?



Example program

$func = o1[name]$

\dots

$o2[name] = func$

\dots

What is name, when $func \mapsto f2$?

\dots

$o2.foo(\dots)$

Critical code example

Analysis state

$o1 = \{foo: f1, bar: f2\}$

$name = T_{str}$

$o2 = \{\}$

$func = f1 | f2$

What is name, when $func \mapsto f1$?

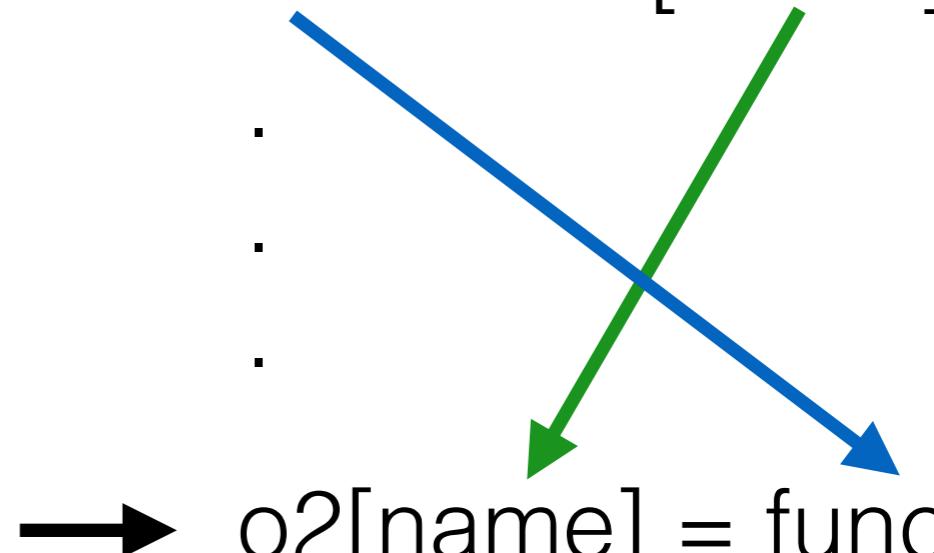
“foo”

What is name, when $func \mapsto f2$?

“bar”

Example program

$func = o1[name]$



Critical code example

Analysis state

$o1 = \{foo: f1, bar: f2\}$

$name = T_{str}$

$o2 = \{foo: f1, bar: f2\}$

$func = f1 | f2$

What is name, when $func \mapsto f1$?

“foo”

What is name, when $func \mapsto f2$?

“bar”

Example program

$func = o1[name]$

$o2[name] = func$

→ $o2.foo(...)$

Critical code example

Analysis state

$o1 = \{foo: f1, bar: f2\}$

$name = T_{str}$

$o2 = \{foo: f1, bar: f2\}$

$func = f1 | f2$

What is name, when $func \mapsto f1$?

“foo”

What is name

“bar”

Example program

$func = o1[name]$

$o2[name] = func$

Resolves only $f1$

→ $o2.foo(...)$

Implementation for JavaScript

- **TAJS_{VR}**: **TAJS** extended with demand-driven value refinement
- **TAJS** is a state-of-the-art analyzer for JavaScript
 - Implemented in Java
 - Active research since 2009
- **VR_{Js}**: Backwards abstract interpreter for JavaScript for answering refinement queries
 - Implemented in Scala from scratch



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Compared to state-of-the-art

	#tests	TAJS	CompAbs	TAJS _{VR}
Underscore ¹	182	0 %	0 %	95% (2.9s)
Lodash3 ¹	176	0 %	0 %	98% (5.5s)
Lodash4 ¹	306	0 %	0 %	87% (24.7s)
Prototype ²	6	0 %	33% (23.1s)	83% (97.7s)
Scriptaculous ²	1	0 %	100% (62.0s)	100% (236.9s)
JQuery ³	71	7% (14.4s)	0 %	7% (17.2s)
JSAI tests ⁴	29	86% (12.3s)	34% (32.4s)	86% (14.3s)

“x% (y)” means succeeded x% of test cases with average time y

¹: Most popular functional utility libraries

²: Wei et al. [2016]

³: Andreasen and Møller [2014]

⁴: Kashyap et al. [2014] & Dewey et al. [2015]

Compared to state-of-the-art

	#tests	TAJS	CompAbs	TAJS _{VR}
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Lodash4 ¹	306	0 %	0 %	87% (24.7s)

Prototype²

6 0 %

33% (23.1s)

Scriptaculous²

1

TAJS_{VR} succeeds analyzing 92% of Underscore and Lodash tests, which all are unanalyzable by existing analyzers

(12.0s)

34% (32.4s)

86% (14.3s)

“x% (y)” means succeeded x% of test cases with average time y

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Value refinement insights

- Value refinement is triggered in few locations
 - In Lodash4, it is triggered in 7 locations in >17000 LoC
- Almost all queries are solved successfully (>99%)
- Queries are answered efficiently (Avg. ~10ms)
- Answering a query requires visiting few locations
 - Typically below 40
- Many queries require interprocedural reasoning

Conclusion

- New technique: Demand-Driven Value Refinement
 - Relational reasoning on top of non-relational analysis
 - Eliminates critical precision loss on-the-fly
 - Uses backwards analysis for gaining relational precision
 - Exploiting forwards analysis state allows efficient refinements
- Experimental evaluation
 - First analysis capable of analyzing most popular JavaScript library
 - No significant overhead for incorporating backwards analyzer
 - Open-source: <https://www.brics.dk/TAJS/VR/>

Value refinement statistics

	Ref locs	Avg # queries	Succ (%)	Refiner time (%)	Avg query time (ms)	Avg. locs visited	Inter (%)
Underscore	5	268	99.98	22.4	2.43	5.05	0.10
Lodash3	12	475	99.99	47.2	5.46	10.47	40.22
Lodash4	7	1284	99.97	52.0	10.01	10.09	25.75
Prototype	4	188	100	2.5	13.08	39.98	48.10
Scriptaculous	2	601	100	3.4	13.21	36.91	42.26
JQuery	5	1	87.5	0.1	13.57	7.1	2.86
JSAI tests	0	-	-	-	-	-	-